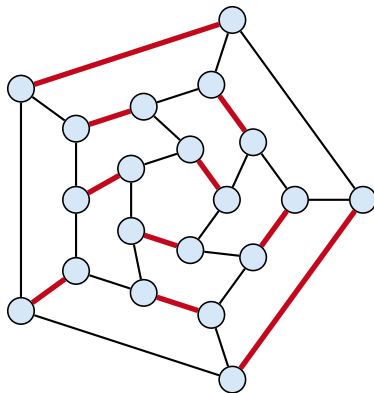


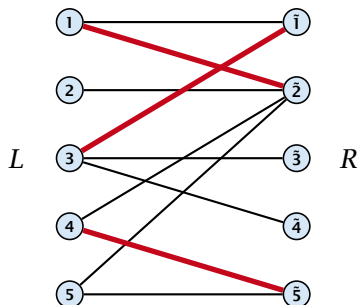
# Matching

- ▶ Input: undirected graph  $G = (V, E)$ .
- ▶  $M \subseteq E$  is a **matching** if each node appears in at most one edge in  $M$ .
- ▶ Maximum Matching: find a matching of maximum cardinality



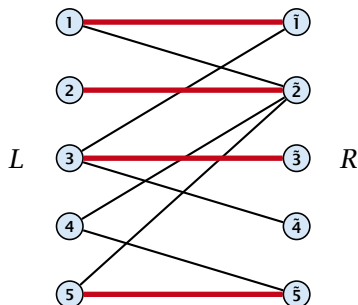
# Bipartite Matching

- ▶ Input: undirected, **bipartite** graph  $G = (L \uplus R, E)$ .
- ▶  $M \subseteq E$  is a **matching** if each node appears in at most one edge in  $M$ .
- ▶ Maximum Matching: find a matching of maximum cardinality



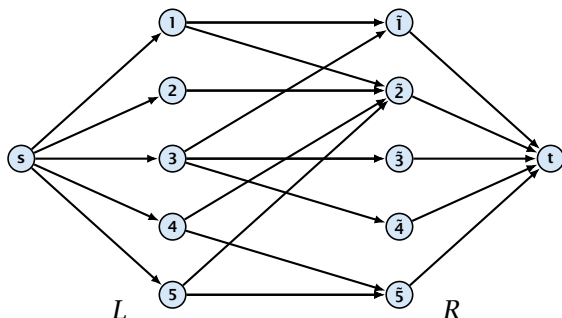
# Bipartite Matching

- ▶ Input: undirected, **bipartite** graph  $G = (L \uplus R, E)$ .
- ▶  $M \subseteq E$  is a **matching** if each node appears in at most one edge in  $M$ .
- ▶ Maximum Matching: find a matching of maximum cardinality



# Maxflow Formulation

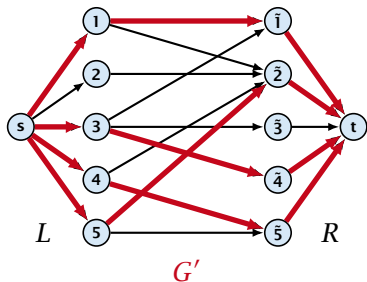
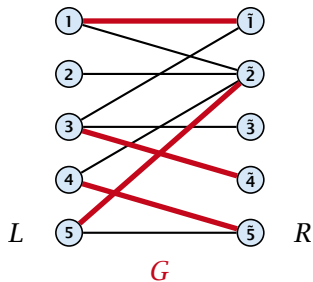
- ▶ Input: undirected, bipartite graph  $G = (L \uplus R \uplus \{s, t\}, E')$ .
- ▶ Direct all edges from  $L$  to  $R$ .
- ▶ Add source  $s$  and connect it to all nodes on the left.
- ▶ Add  $t$  and connect all nodes on the right to  $t$ .
- ▶ All edges have unit capacity.



# Proof

## Max cardinality matching in $G \leq$ value of maxflow in $G'$

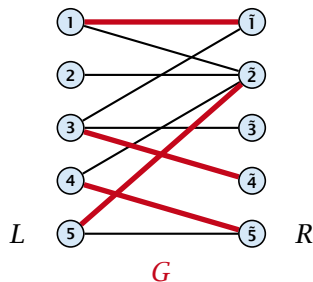
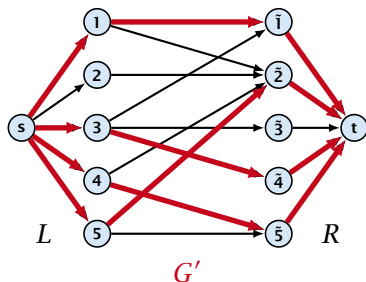
- ▶ Given a maximum matching  $M$  of cardinality  $k$ .
- ▶ Consider flow  $f$  that sends one unit along each of  $k$  paths.
- ▶  $f$  is a flow and has cardinality  $k$ .



# Proof

Max cardinality matching in  $G \geq$  value of maxflow in  $G'$

- ▶ Let  $f$  be a maxflow in  $G'$  of value  $k$
- ▶ Integrality theorem  $\Rightarrow k$  integral; we can assume  $f$  is 0/1.
- ▶ Consider  $M =$  set of edges from  $L$  to  $R$  with  $f(e) = 1$ .
- ▶ Each node in  $L$  and  $R$  participates in at most one edge in  $M$ .
- ▶  $|M| = k$ , as the flow must use at least  $k$  middle edges.



# 8.1 Matching

## Which flow algorithm to use?

- ▶ Generic augmenting path:  $\mathcal{O}(m \text{val}(f^*)) = \mathcal{O}(mn)$ .
- ▶ Capacity scaling:  $\mathcal{O}(m^2 \log C) = \mathcal{O}(m^2)$ .
- ▶ Shortest augmenting path:  $\mathcal{O}(mn^2)$ .

For **unit capacity simple graphs** shortest augmenting path can be implemented in time  $\mathcal{O}(m\sqrt{n})$ .

A graph is a **unit capacity simple graph** if

- ▶ every edge has capacity 1
- ▶ a node has either at most one leaving edge **or** at most one entering edge

# Baseball Elimination

<i>team</i> <i>i</i>	<i>wins</i> $w_i$	<i>losses</i> $\ell_i$	<i>remaining games</i>			
			<i>Atl</i>	<i>Phi</i>	<i>NY</i>	<i>Mon</i>
Atlanta	83	71	–	1	6	1
Philadelphia	80	79	1	–	0	2
New York	78	78	6	0	–	0
Montreal	77	82	1	2	0	–

**Which team can end the season with most wins?**

- ▶ Montreal is eliminated, since even after winning all remaining games there are only 80 wins.
- ▶ But also Philadelphia is eliminated. Why?



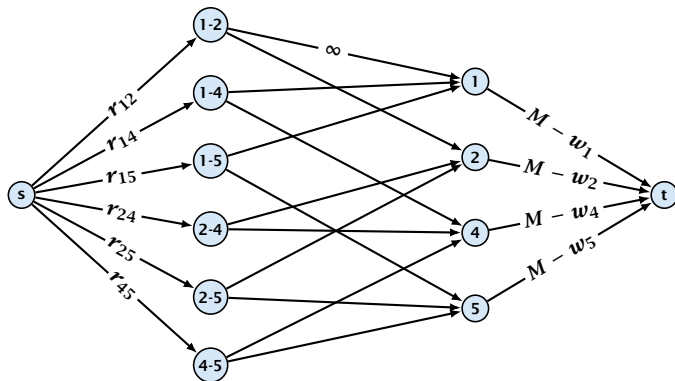
# Baseball Elimination

## Formal definition of the problem:

- ▶ Given a set  $S$  of teams, and one specific team  $z \in S$ .
- ▶ Team  $x$  has already won  $w_x$  games.
- ▶ Team  $x$  still has to play team  $y$ ,  $r_{xy}$  times.
- ▶ Does team  $z$  still have a chance to finish with the most number of wins.

# Baseball Elimination

Flow network for  $z = 3$ .  $M$  is number of wins Team 3 can still obtain.




**Idea.** Distribute the results of remaining games in such a way that no team gets too many wins.

# Certificate of Elimination

Let  $T \subseteq S$  be a subset of teams. Define

$$w(T) := \sum_{i \in T} w_i, \quad r(T) := \sum_{i, j \in T, i < j} r_{ij}$$



If  $\frac{w(T)+r(T)}{|T|} > M$  then one of the teams in  $T$  will have more than  $M$  wins in the end. A team that can win at most  $M$  games is therefore eliminated.

## Theorem 51

A team  $z$  is eliminated if and only if the flow network for  $z$  does not allow a flow of value  $\sum_{i \in S \setminus \{z\}, i < j} r_{ij}$ .

### Proof ( $\Leftarrow$ )

- ▶ Consider the mincut  $A$  in the flow network. Let  $T$  be the set of team-nodes in  $A$ .
- ▶ If for node  $x-y$  not both team-nodes  $x$  and  $y$  are in  $T$ , then  $x-y \notin A$  as otw. the cut would cut an infinite capacity edge.
- ▶ We don't find a flow that saturates all source edges:

$$\begin{aligned}r(S \setminus \{z\}) &> \text{cap}(A, V \setminus A) \\ &\geq \sum_{i < j: i \notin T \vee j \notin T} r_{ij} + \sum_{i \in T} (M - w_i) \\ &\geq r(S \setminus \{z\}) - r(T) + |T|M - w(T)\end{aligned}$$

- ▶ This gives  $M < (w(T) + r(T))/|T|$ , i.e.,  $z$  is eliminated.

# Baseball Elimination

## Proof ( $\Rightarrow$ )

- ▶ Suppose we have a flow that saturates all source edges.
- ▶ We can assume that this flow is **integral**.
- ▶ For every pairing  $x$ - $y$  it defines how many games team  $x$  and team  $y$  should win.
- ▶ The flow leaving the team-node  $x$  can be interpreted as the additional number of wins that team  $x$  will obtain.
- ▶ This is less than  $M - w_x$  because of capacity constraints.
- ▶ Hence, we found a set of results for the remaining games, such that no team obtains more than  $M$  wins in total.
- ▶ Hence, team  $z$  is not eliminated.

# Project Selection

## Project selection problem:

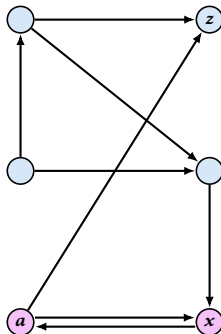
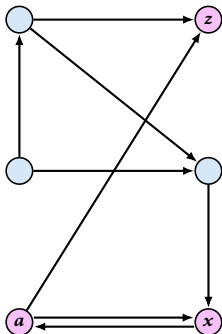
- ▶ Set  $P$  of possible projects. Project  $v$  has an associated profit  $p_v$  (can be positive or negative).
- ▶ Some projects have requirements (taking course EA2 requires course EA1).
- ▶ Dependencies are modelled in a graph. Edge  $(u, v)$  means “can’t do project  $u$  without also doing project  $v$ .”
- ▶ A subset  $A$  of projects is **feasible** if the prerequisites of every project in  $A$  also belong to  $A$ .

**Goal:** Find a feasible set of projects that maximizes the profit.

# Project Selection

## The prerequisite graph:

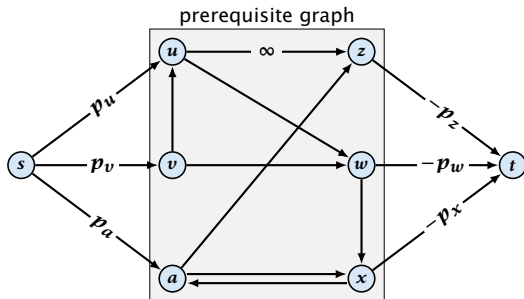
- ▶  $\{x, a, z\}$  is a feasible subset.
- ▶  $\{x, a\}$  is infeasible.



# Project Selection

## Mincut formulation:

- ▶ Edges in the prerequisite graph get infinite capacity.
- ▶ Add edge  $(s, v)$  with capacity  $p_v$  for nodes  $v$  with positive profit.
- ▶ Create edge  $(v, t)$  with capacity  $-p_v$  for nodes  $v$  with negative profit.





## Theorem 52

$A$  is a mincut if  $A \setminus \{s\}$  is the optimal set of projects.

### Proof.

▶  $A$  is feasible because of capacity infinity edges.

$$\text{cap}(A, V \setminus A) = \sum_{v \in \bar{A}: p_v > 0} p_v + \sum_{v \in A: p_v < 0} (-p_v)$$

$$= \sum_{v: p_v > 0} p_v - \sum_{v \in A} p_v$$

For the formula we define  $p_s := 0$ .

The step follows by adding  $\sum_{v \in A: p_v > 0} p_v - \sum_{v \in A: p_v > 0} p_v = 0$ .

Note that minimizing the capacity of the cut  $(A, V \setminus A)$  corresponds to maximizing profits of projects in  $A$ .

