8 Seidels LP-algorithm

- Suppose we want to solve $\min\{c^Tx \mid Ax \geq b; x \geq 0\}$, where $x \in \mathbb{R}^d$ and we have m constraints.
- In the worst-case Simplex runs in time roughly $\mathcal{O}(m(m+d)\binom{m+d}{m}) \approx (m+d)^m$. (slightly better bounds on the running time exist, but will not be discussed here).
- If d is much smaller than m one can do a lot better.
- In the following we develop an algorithm with running time $\mathcal{O}(d! \cdot m)$, i.e., linear in m.

8 Seidels LP-algorithm

Setting:

- We assume an LP of the form

$$\begin{array}{lrcl} \min & c^Tx & & \\ \text{s.t.} & Ax & \geq & b \\ & x & \geq & 0 \end{array}$$

- We assume that the LP is bounded.

Ensuring Conditions

Given a standard minimization LP

$$\begin{array}{lrcl} \min & c^Tx & & \\ \text{s.t.} & Ax & \geq & b \\ & x & \geq & 0 \end{array}$$

how can we obtain an LP of the required form?

- **Compute a lower bound on c^Tx for any basic feasible solution.**

Computing a Lower Bound

Let s denote the smallest common multiple of all denominators of entries in A, b.

Multiply entries in A, b by s to obtain integral entries. This does not change the feasible region.

Add slack variables to A; denote the resulting matrix with $\bar{A}$.

If B is an optimal basis then x_B with $\bar{A}_B x_B = \bar{b}$, gives an optimal assignment to the basis variables (non-basic variables are 0).

Theorem 2 (Cramers Rule)
Let M be a matrix with $\det(M) \neq 0$. Then the solution to the system $Mx = b$ is given by

$$x_i = \frac{\det(M_j)}{\det(M)} ,$$

where M_i is the matrix obtained from M by replacing the i-th column by the vector b.

Proof:

- Define
$$X_i = \begin{pmatrix} | & & | & | & | & & | \\ e_1 & \cdots & e_{i-1} & x & e_{i+1} & \cdots & e_n \\ | & & | & | & | & & | \end{pmatrix}$$
Note that expanding along the i-th column gives that $\det(X_i) = x_i$.
- Further, we have
$$MX_i = \begin{pmatrix} | & & | & | & | & & | \\ Me_1 & \cdots & Me_{i-1} & Mx & Me_{i+1} & \cdots & Me_n \\ | & & | & | & | & & | \end{pmatrix} = M_i$$
- Hence,
$$x_i = \det(X_i) = \frac{\det(M_i)}{\det(M)}$$

Bounding the Determinant

Let Z be the maximum absolute entry occuring in $\bar{A}, \bar{b}$ or c. Let C denote the matrix obtained from $\bar{A}_B$ by replacing the j-th column with vector $\bar{b}$ (for some j).

Observe that

$$\begin{aligned} |\det(C)| &= \left| \sum_{\pi \in S_m} \mathrm{sgn}(\pi) \prod_{1 \le i \le m} C_{i\pi(i)} \right| \\ &\le \sum_{\pi \in S_m} \prod_{1 \le i \le m} |C_{i\pi(i)}| \\ &\le m! \cdot Z^m . \end{aligned}$$

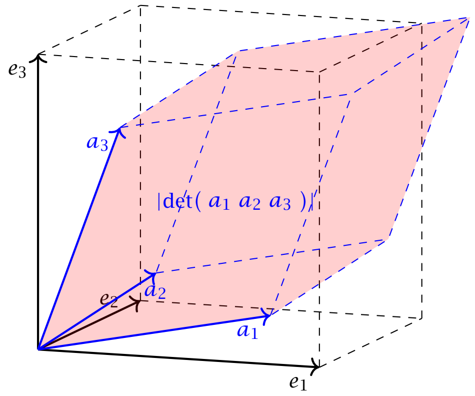
Here $\mathrm{sgn}(\pi)$ denotes the sign of the permutation, which is 1 if the permutation can be generated by an even number of transpositions (exchanging two elements), and -1 if the number of transpositions is odd.
The first identity is known as Leibniz formula.

Bounding the Determinant

Alternatively, Hadamards inequality gives

$$\begin{aligned} |\det(C)| &\le \prod_{i=1}^{m} \|C_{*i}\| \le \prod_{i=1}^{m} (\sqrt{m} Z) \\ &\le m^{m/2} Z^m . \end{aligned}$$

Hadamards Inequality



Hadamards inequality says that the volume of the red parallelepiped (Spat) is smaller than the volume in the black cube (if $\|e_1\| = \|a_1\|$, $\|e_2\| = \|a_2\|$, $\|e_3\| = \|a_3\|$).

Ensuring Conditions

Given a standard minimization LP

$$\begin{array}{llll} \min & c^T x & & \\ \text{s.t.} & Ax & \geq & b \\ & x & \geq & 0 \end{array}$$

how can we obtain an LP of the required form?

- **Compute a lower bound on $c^T x$ for any basic feasible solution.** Add the constraint $c^T x \geq -dZ(m! \cdot Z^m) - 1$. Note that this constraint is superfluous unless the LP is unbounded.

Ensuring Conditions

Compute an optimum basis for the new LP.

- If the cost is $c^T x = -(dZ)(m! \cdot Z^m) - 1$ we know that the original LP is unbounded.
- Otw. we have an optimum basis.

In the following we use $\mathcal{H}$ to denote the set of all constraints apart from the constraint $c^T x \geq -dZ(m! \cdot Z^m) - 1$.

We give a routine SeidelLP$(\mathcal{H}, d)$ that is given a set $\mathcal{H}$ of explicit, non-degenerate constraints over d variables, and minimizes $c^T x$ over all feasible points.

In addition it obeys the implicit constraint $c^T x \geq -(dZ)(m! \cdot Z^m) - 1$.

Algorithm 1 SeidelLP($\mathcal{H}, d$)

```
if d = 1 then solve 1-dimensional problem and return;
if H = ∅ then return x on implicit constraint hyperplane
choose random constraint h ∈ H
Ĥ ← H \ {h}
x̂* ← SeidelLP(Ĥ, d)
if x̂* = infeasible then return infeasible
if x̂* fulfills h then return x̂*
// optimal solution fulfills h with equality, i.e., a_h^T x = b_h
solve a_h^T x = b_h for some variable x_ℓ;
eliminate x_ℓ in constraints from Ĥ and in implicit constr.;
x̂* ← SeidelLP(Ĥ, d − 1)
if x̂* = infeasible then
      return infeasible
else
      add the value of x_ℓ to x̂* and return the solution
```

8 Seidels LP-algorithm

Note that for the case $d = 1$, the asymptotic bound $\mathcal{O}(\max\{m, 1\})$ is valid also for the case $m = 0$.

- If $d = 1$ we can solve the 1-dimensional problem in time $\mathcal{O}(\max\{m, 1\})$.
- If $d > 1$ and $m = 0$ we take time $\mathcal{O}(d)$ to return d-dimensional vector x.
- The first recursive call takes time $T(m-1, d)$ for the call plus $\mathcal{O}(d)$ for checking whether the solution fulfills h.
- If we are unlucky and $\hat{x}^*$ does not fulfill h we need time $\mathcal{O}(d(m+1)) = \mathcal{O}(dm)$ to eliminate x_ℓ. Then we make a recursive call that takes time $T(m-1, d-1)$.
- The probability of being unlucky is at most d/m as there are at most d constraints whose removal will decrease the objective function

8 Seidels LP-algorithm

This gives the recurrence

$$T(m,d) = \begin{cases} \mathcal{O}(\max\{1, m\}) & \text{if } d = 1 \\ \mathcal{O}(d) & \text{if } d > 1 \text{ and } m = 0 \\ \mathcal{O}(d) + T(m-1, d) + & \\ \frac{d}{m}(\mathcal{O}(dm) + T(m-1, d-1)) & \text{otw.} \end{cases}$$

Note that $T(m, d)$ denotes the expected running time.

8 Seidels LP-algorithm

Let C be the largest constant in the $\mathcal{O}$-notations.

$$T(m,d) = \begin{cases} C\max\{1, m\} & \text{if } d = 1 \\ Cd & \text{if } d > 1 \text{ and } m = 0 \\ Cd + T(m-1, d) + & \\ \frac{d}{m}(Cdm + T(m-1, d-1)) & \text{otw.} \end{cases}$$

Note that $T(m, d)$ denotes the expected running time.

8 Seidels LP-algorithm

Let C be the largest constant in the $\mathcal{O}$-notations.

We show $T(m,d) \le Cf(d)\max\{1,m\}$.

$d = 1$:

$$T(m,1) \le C\max\{1,m\} \le Cf(1)\max\{1,m\} \text{ for } f(1) \ge 1$$

$d > 1; m = 0$:

$$T(0,d) \le \mathcal{O}(d) \le Cd \le Cf(d)\max\{1,m\} \text{ for } f(d) \ge d$$

$d > 1; m = 1$:

$$\begin{aligned} T(1,d) &= \mathcal{O}(d) + T(0,d) + d\left(\mathcal{O}(d) + T(0,d-1)\right) \\ &\le Cd + Cd + Cd^2 + dCf(d-1) \\ &\le Cf(d)\max\{1,m\} \text{ for } f(d) \ge 3d^2 + df(d-1) \end{aligned}$$

8 Seidels LP-algorithm

$d > 1; m > 1$:
(by induction hypothesis statm. true for $d' < d, m' \ge 0$;
and for $d' = d$, $m' < m$)

$$\begin{aligned} T(m,d) &= \mathcal{O}(d) + T(m-1,d) + \frac{d}{m}\left(\mathcal{O}(dm) + T(m-1,d-1)\right) \\ &\le Cd + Cf(d)(m-1) + Cd^2 + \frac{d}{m}Cf(d-1)(m-1) \\ &\le 2Cd^2 + Cf(d)(m-1) + dCf(d-1) \\ &\le Cf(d)m \end{aligned}$$

if $f(d) \ge df(d-1) + 2d^2$.

8 Seidels LP-algorithm

- Define $f(1) = 3 \cdot 1^2$ and $f(d) = df(d-1) + 3d^2$ for $d > 1$.

Then

$$\begin{aligned} f(d) &= 3d^2 + df(d-1) \\ &= 3d^2 + d\left[3(d-1)^2 + (d-1)f(d-2)\right] \\ &= 3d^2 + d\left[3(d-1)^2 + (d-1)\left[3(d-2)^2 + (d-2)f(d-3)\right]\right] \\ &= 3d^2 + 3d(d-1)^2 + 3d(d-1)(d-2)^2 + \dots \\ &\quad + 3d(d-1)(d-2)\cdot\ldots\cdot 4\cdot 3\cdot 2\cdot 1^2 \\ &= 3d!\left(\frac{d^2}{d!} + \frac{(d-1)^2}{(d-1)!} + \frac{(d-2)^2}{(d-2)!} + \dots\right) \\ &= \mathcal{O}(d!) \end{aligned}$$

since $\sum_{i\ge1}\frac{i^2}{i!}$ is a constant.

$$\sum_{i\ge1}\frac{i^2}{i!} = \sum_{i\ge0}\frac{i+1}{i!} = e + \sum_{i\ge1}\frac{i}{i!} = 2e$$